I.MX MEMORY MADNESS

HOW TO DUMP, PARSE, AND ANALYZE I.MX FLASH MEMORY CHIPS

Damien Cauquil | HITB Amsterdam 2019 (🖄 🔌)

WHO AM I?

- **Q** Head of R&D @ Econocom Digital.Security
- **T** Senior security researcher
- Hardware hacker (or at least pretending)

AGENDA

- Firmware extraction 101
- Meet the i.MX architecture
- i.MX flash memory layout
- imx-nand-tools FTW
- Best practices

FIRMWARE EXTRACTION 101

WHY DO WE WANT TO EXTRACT A DEVICE'S FIRMWARE ?

- Contains filesystems, applications, binary files
- May also contain **VERY interesting data**: encryption/decryption keys, certificates, passwords

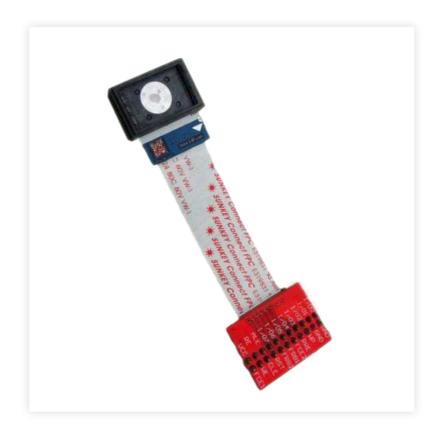
WHY DO WE WANT TO EXTRACT A DEVICE'S FIRMWARE ?

<u>We need to understand everything about a device:</u>

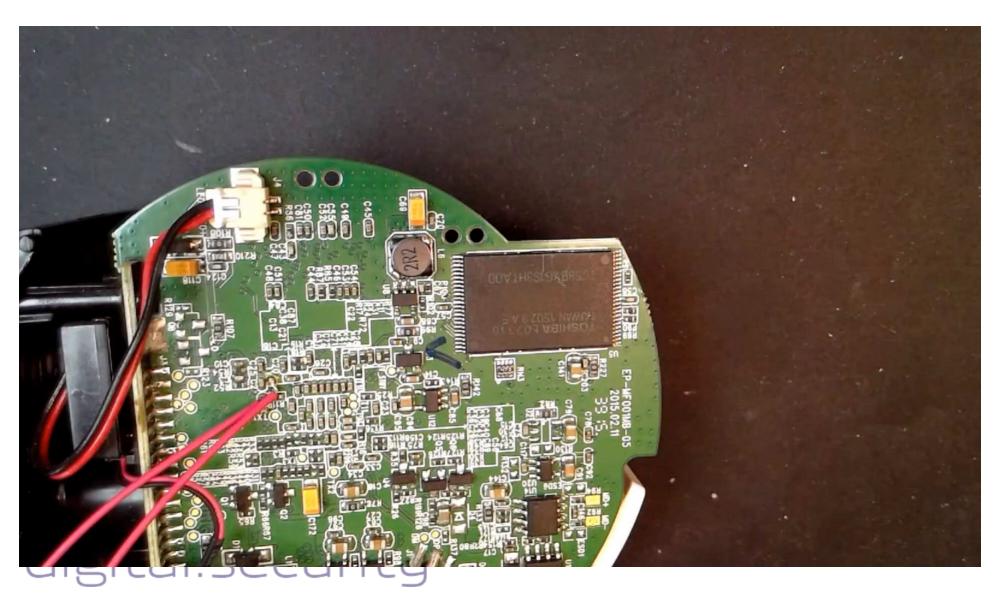
- How it has been designed
- How it (really) works

• Where and how every bit of data is stored digital.security

METHOD #1: CLIPPING & READING



METHOD #2: CHIP-OFF

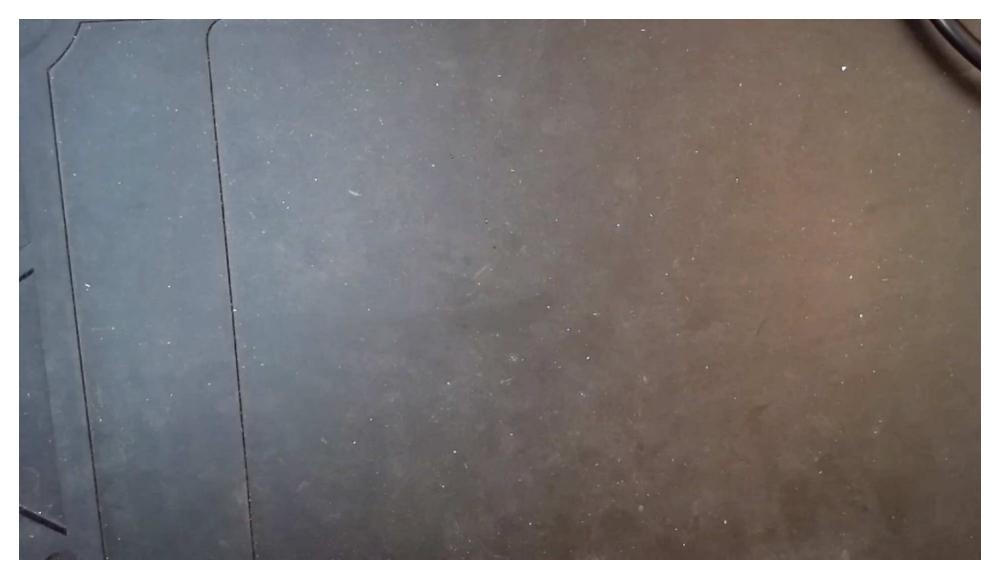


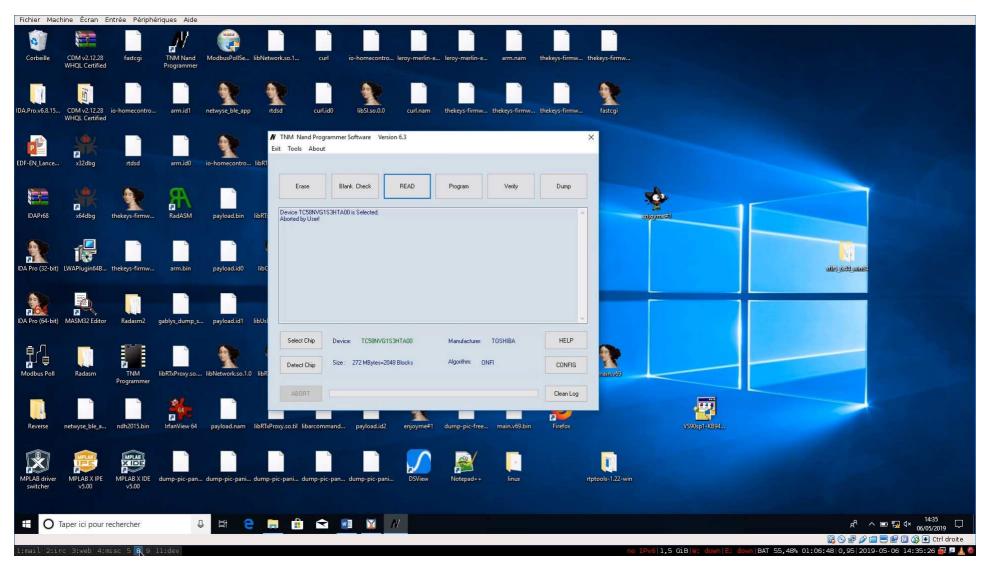
PROFESSIONAL FLASH PROGRAMMER











NAND DUMP SIZE

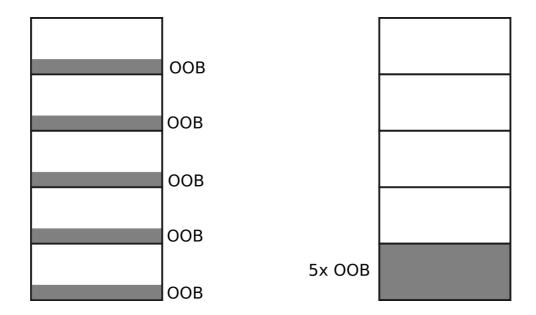
\$ ls -alh camera.bin
-rwx----- 1 virtualabs virtualabs 1,1G camera.bin

Dump file is greater than 1 GB !

PAGES, BYTES AND OOB

- Bytes are stored, erased, and modified in pages
- NAND flash chips are not 100% reliable and errors when storing bits may occur
- To avoid this, vendors usually provide more space to store Error Correction Codes (ECC) in spare-byte area (OOB) digital.security

PAGES, BYTES AND OOB



PAGES, BYTES AND OOB

MT29F16G08ADBCA, MT29F16G16ADBCA

Features

- Open NAND Flash Interface (ONFI) 1.0-compliant¹
- Single-level cell (SLC) technology

Organization

- Page size x8: 4320 bytes (4096 + 224 bytes)
- Page size x16: 2160 words (2048 + 112 words)
- Block size: 64 pages (256K + 14K bytes)
- Plane size: 2 planes x 2048 blocks per plane
- Device size: 8Gb: 4096 blocks
- Device size: 16Gb: 8192 blocks
- Asynchronous I/O performance
 - ^tRC/^tWC: 20ns (3.3V), 30ns (1.8V)

- First l
 - ped fi ECC,
- RESE
 - powe
- Altern
 powe
- Interi
 - plane
- Quali
 - Da
 - cat
- En

REMOVING THE OOB DATA

Hex Edit - [camera.bin]	-
<u> </u>	
🖙 🖬 🚭 🖪 🐰 🖻 🛍 🗠 😭 🗛 🏹 🤐 🏭 - 🎌 - 🎇 - 👫 - ASCII default 🛛 🛛 🕑 🥒 🔹 🕨	
🙀 🗸 🗸 🙀 🕺 🗸 🖓 🐜 🖓 🖓 🖓 🖓	
LEROY-~2.BIN Camera.bin	
027D 11B00: E3 00 00 97 CD 00 97 CD 00 AD E1 02 00 M. L0 M. 027D 11C0: 13 E3 12 5E 06 08 00 02 E5 01 60 AO E1 02 00 M. L0 027D 11D0: 13 E3 12 5E 06 08 00 02 10 03 56 0F GF	

CHECKING OUR DUMP WITH BINWALK

<pre>\$ binwalk</pre>	ipcam.fw.bin	
DECIMAL	HEXADECIMAL	DESCRIPTION
96188 []	0x177BC	CRC32 polynomial table, []
2490368	0x260000	Squashfs filesystem, []
4456448	0x440000	Squashfs filesystem, []
5505024	0x540000	Squashfs filesystem, []
6684672	0x660000	Squashfs filesystem, []
7208960	0x6E0000	JFFS2 filesystem, little endian
7643512	0x74A178	JFFS2 filesystem, little endian

EXTRACTING FILES FROM VARIOUS FILESYSTEMS

- SquashFS: compressed filesystem, one partition/image
- YAFFS2: Yet Another Flash FS
- JFFS2: Journalized Flash FS version 2, one partition/image
- UBI: Unsorted Block Image, multiple partitions with various FS digital.security



IT'S A DOCUMENTED PROCESS

PenTestPartners just published a blog entry:

http://bit.ly/HITB-PTPFW

AND WE STUMBLED UPON AN I.MX6 SYSTEM



digital.security

I

virtualabs@virtubox:~\$



HEX ANALYSIS REVEALED WEIRD BYTES

3D81:E2E0 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E2F0 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E300 00000000	00 00 00 00	000000000	00 00 00 00	
3D81:E310 00000000	000000000	000000000	00 00 00 00	
3D81:E320 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E330 <u>00</u> 000000	00 00 00 00	000000000	00 00 00 00	
3D81:E340 FF 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	ÿ
3D81:E350 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E360 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E370 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E380 00000000	000000000	00 00 00 00	00 00 00 00	<u>.</u>
3D81:E390 00000000	00 00 00 00	00 00 00	00 00 00 00	
3D81:E3A0 00000000	000000000	00 00 00 00	00 00 00 00	
3D81:E3B0 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E3C0 00000000	000000000	00 00 00 00	00 00 00 00	
3D81:E3D0 00000000	00 00 00 00	000000000	00 00 00 00	
3D81:E3E0 00000000	00 00 00 00	00 00 00 00	00 00 00 00	

A CRAPPY BYTE BEFORE UBI SIGNATURE

<i>YYYYYYYYYYYYYY YYYYYYYYYY VVVVV ~~~~~~ YYYYYYYYYYYYYYYYY **30C3**: **4800** 00 **5**5 42 49 23 01 00 00 00 00 00 00 000000000 UBT#. 30C3:4810 020000 10 000000 20 0007 42 E4 66 00 00 00Bäf 30C3:4830 00000000 0000000 00 15 87 C5 00 00 00 00 30C3:4840 78000000 00000000 00 00 00 00 00 00 00 00 Χ. 30C3:4850 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00

SAME 1-BYTE OFFSET IN BINWALK OUTPUT

24891865	0x17BD1D9	YAFFS filesystem
25159701	0x17FE815	YAFFS filesystem
25436181	0x1842015	YAFFS filesystem
25712661	0x1885815	YAFFS filesystem
25727234	0x1889102	Unix path: /usr/share/brw/local/index.html 0
35389441	0x21C0001	UBI erase count header, version: 1, EC: 0x4, VID header offset: 0x1000

UBI header is not aligned on page size (0x1000)

THAT'S WEIRD 😳

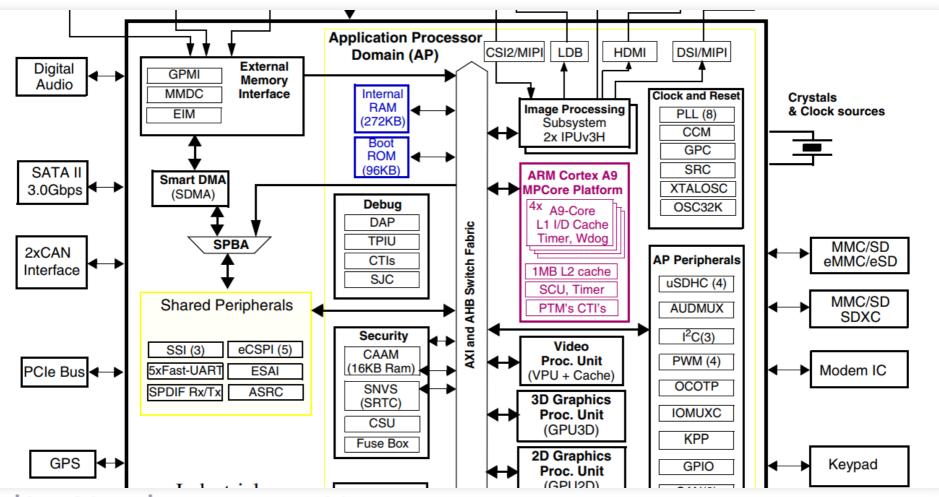
- Quick investigation revealed **anomalies**
- Our **dump seems OK**, but we still cannot extract data from it
- It must be related to **i.MX**: maybe a custom storage mechanism

I.MX ARCHITECTURE AND MEMORY LAYOUT

I.MX ARCHITECTURE

- Integrated Multimedia Application processors
- Popular in automotive and home automation industries
- Provides a lot of features including:
 - Secure/non-secure RAM
 - SATA II support
- digsears Boot mity

I.MX ARCHITECTURE



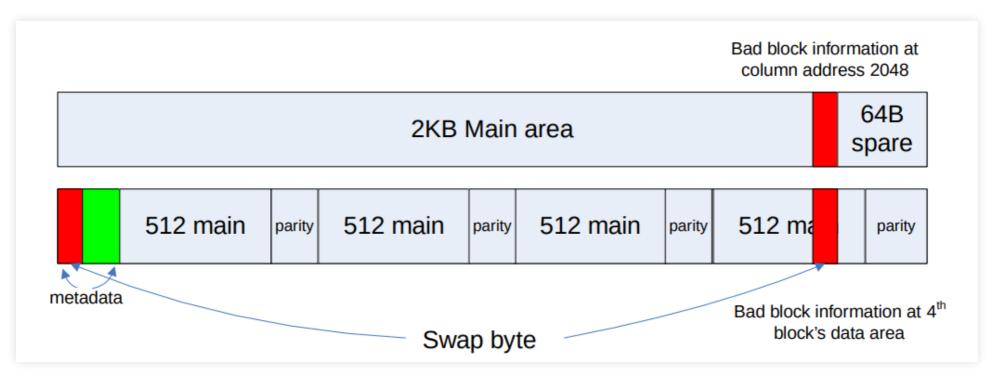
I.MX ARCHITECTURE

- Can boot on various storage devices:
 - NAND Flash
 - Parallel NOR Flash
 - SD card
 - MMC
 - SATA HDD
- It also embeds a boot ROM (Freescale Inc.) digital.security

GENERAL-PURPOSE MULTIMEDIA INTERFACE

- controls how data is read/stored on NAND flash chips
- supports multiple NAND flash chips
- uses **BCH** to perform error control and correction digital.security

NAND FLASH STRUCTURE



(image extracted from i.MX28 reference manual)

HOW IS DATA STORED ?

- Data is split in **512-byte chunks**
- ECC bits are added at the end of each chunk
- Chunks are then **grouped and stored in a page** preceeded by one metadata block
- Bad block marker byte is swapped with first

WEIRD BYTE EXPLAINED !

3D81:E2E0 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E2F0 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E300 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E310 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E320 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E330 00 00 00 00			00 00 00 00	
3D81:E340 FF 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	ÿ
3D81:E350 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E360 00000000	00 00 00 00	00 00 00 00	000000000	
3D81:E370 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E380 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E390 00000000	00 00 00 00	00 00 00	00 00 00 00	
3D81:E3A0 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E3B0 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E3C0 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E3D0 00000000	00 00 00 00	00 00 00 00	00 00 00 00	
3D81:E3E0 00000000	00 00 00 00	00 00 00 00	00 00 00 00	

FIRMWARE CONFIGURATION BLOCK (FCB)

- This structure contains **all the required information** about how data is stored
- It must be present in the **first 1MB**
- Second field of this structure contains "FCB " in ASCII

FCB SIGNATURE IN HEXDUMP

0000:0000 00 00 00 00	00 00 00 00	00 00 00 00	OD FB FF FF	<u></u> ûÿÿ
0000:0010 46 43 42 20	00 00 00 01	50 3C 19 06	00 00 00 00	FCBP<
0000:0020 00 10 00 00	E0 10 00 00	40 00 00 00	00 00 00 00	à@
0000:0030 00 00 00 00	00 00 00 00	08 00 00 00	00 02 00 00	
0000:0040 00 02 00 00	08 00 00 00	01 00 00 00	07 00 00 00	
0000:0050 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0000:0060 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0000:0070 00 00 00 00	C2 15 00 00	42 11 00 00	36 00 00 00	ÄB6
0000:0080 36 00 00 00	00 01 00 00	49 OF 00 00	00 00 00 00	6I
0000:0090 00100000	00 00 00 00	00 00 00 00	00 00 00 00	
0000:00A0 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0000:00B0 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0000:0000 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0000:00D0 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0000:00E0 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0000:00F0 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0000:0100 00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	

FIRMWARE CONFIGURATION BLOCK (FCB)

- NAND page data size
- Block N ECC type
- Block N size
- Block 0 ECC type
- Block 0 size
- Number of bytes in metadata of a page

FCB SIGNATURE IN HEXDUMP



Offset +0x3C: number of bytes of metadata block digital.security

1-BYTE OFFSET EXPLAINED!

<i>YYYYYYYYYYYYY **YYYYYYYYYYYY** VVVVV ~~~~~~ YYYYYYYYYYYYYYYYY **30C3**: **4800** 00 **5**5 42 49 23 01 00 00 00 00 00 00 UBT# 00 00 00 00 30C3:4810 02 00 00 10 00 00 00 20 00 07 42 E4 66 00 00 00Bäf 30C3:4830 00 00 00 00 00 00 00 00 00 00 15 87 C5 00 00 00 00 30C3:4840 78000000 00000000 00 00 00 00 00 00 00 00 Χ. 30C3:4850 00 00 00 00 00 00 00 00 00

DISCOVERED BAD BLOCK TABLE (DBBT)

- Provides custom NAND bad block management
- Its headers provide information about the number of bad blocks and impacted pages

0010:E000	00 00 00 00	00 44 42 42	54 00 00 00	01 00 00 00	DBBT
0010:E010	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0010:E020	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0010:E030	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0010:E040	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0010:E050	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0010:E060	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0010:E070	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	
0010:E080	00 00 00 00	00 00 00 00	00 00 00 00	00 00 00 00	

ECC

M 512 bytes EccB 0 512 bytes EccBN 512 bytes EccBN 512 bytes EccBN	ſ		Block 0		Block 1		Block 2		Block 3		
		М		EccB 0		EccBN		EccBN		EccBN	

(image extracted from i.MX28 reference manual)

ECC

- Provides a way to dynamically fix errors, if possible
- Uses **BCH** (Bose, Ray-Chaudhuri and Hocquenghem) error-correcting code
- Data bytes may be shifted by a number of bits due to BCH bits

SO, WHAT'S NEXT?

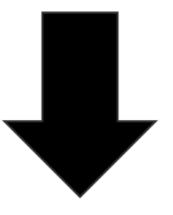
FROM NAND FLASH DUMP TO FILESYSTEMS

RECOVER AND REMAP ALL THE BYTES

- We first find an **FCB structure** and parse it to recover all the critical parameters
- Then we **remove every metadata and ECC bits** according to this FCB
- We use ECC bits to fix errors and save each block in an output file digital.security

@ Page address

META DATA Block 0	ECC	Block 1	ECC	Block 2		ECC
----------------------	-----	---------	-----	---------	--	-----



@ rectified page address

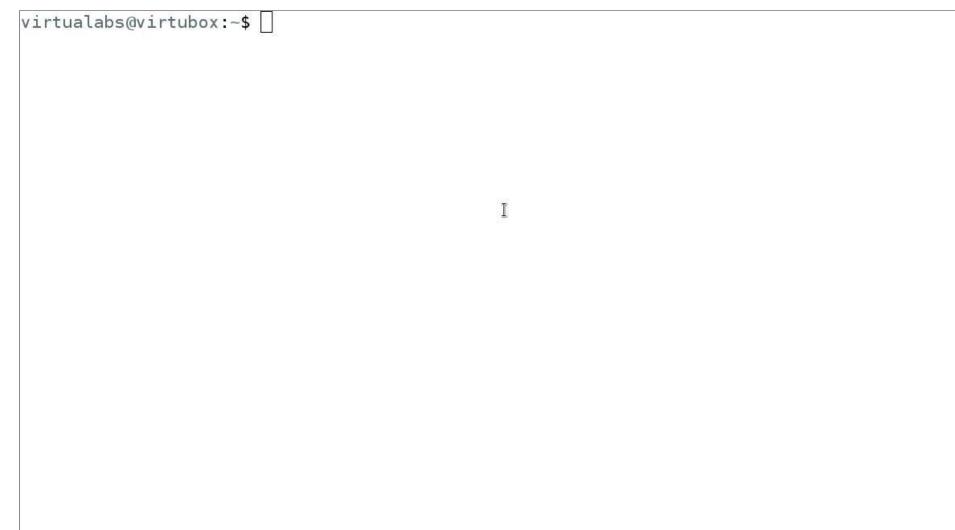
Block 0	Block 1	Block 2	
---------	---------	---------	--

IMX NAND TOOLS

\$ sudo pip install imx-nand-tools

https://github.com/DigitalSecurity/imx-nand-tools/

FCB PARSING



algital.security

CONVERTING IMAGE TO USEABLE DUMP

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algrean.seearreg

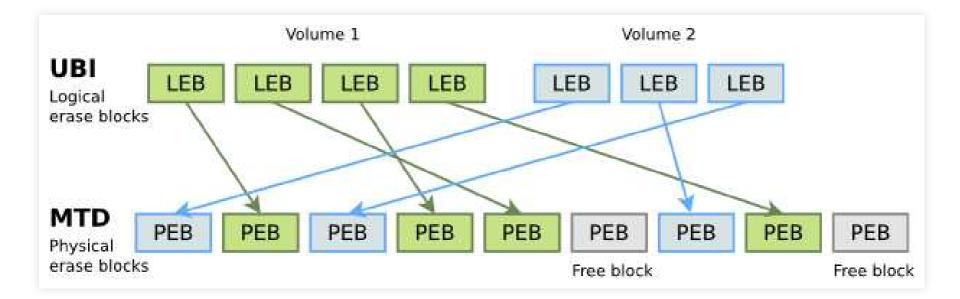
ANALYZING THIS NEW DUMP

I

virtualabs@virtubox:~\$ 🗌

agra.secunty

UBI OVERVIEW



UBIREADER

- Provides a **set of tools** to parse, analyze and extract volumes and files from a UBI container
- **Open-source** and available on Github
- Written in **Python**
- Does not support *fastboot* mode digital.security

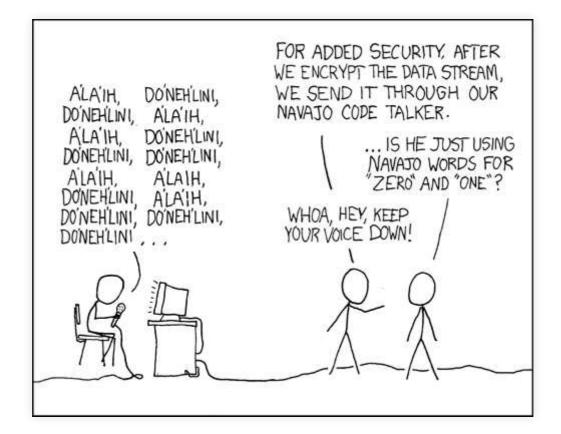
ACCESSING FILES STORED IN VARIOUS IMAGES

```
$ ubireader extract files -iw img-xx vol-iio 0633 0.ubifs
[...]
$ ls ubifs-root/ -al
total 76
drwxr-xr-x 19 virtualabs virtualabs 4096 mai
                                                 9 09:40
drwxr-xr-x 3 virtualabs virtualabs 4096 mai
                                                 9 09:40
                                                 9 09:40
drwxr-xr-x 2 virtualabs virtualabs 4096 mai
                                                        bin
drwxr-xr-x 2 virtualabs virtualabs 4096 mai
                                                 9 09:40 hoot
drwxr-xr-x 5 virtualabs virtualabs 4096 mai
                                                 9 09:40 Data
[\ldots]
drwxr-xr-x
            2 virtualabs virtualabs 4096 mai
                                                 9 09:40 tmp
            7 virtualabs virtualabs 4096 mai
                                                 9 09:40 usr
drwxr-xr-x
            2 virtualabs virtualabs 4096 mai
                                                 9 09:40 var
drwxr-xr-x
```

THAT'S A WIN



SECURITY THROUGH OBSCURITY



(Image: XKCD #257)

NOT SO OBSCURE AFTERALL

- **Reference manuals** describe how i.MX GPMI works and how data is read/stored on NAND flash memory
- **Publicly available code on Github** provides a better understanding of critical structures and how things are implemented

IMX KNOBS GITHUB REPOSITORY

struct fcb_block {

F	CB_ROM_NAND_Tim	ning_t m_NANDTiming;	//!< Optimum timing parameters for Tas, Tds, Tdh in nsec.	
u	int32_t	m_u32PageDataSize;	//!< 2048 for 2K pages, 4096 for 4K pages.	
u	int32_t	m_u32TotalPageSize;	//!< 2112 for 2K pages, 4314 for 4K pages.	
u	int32_t	<pre>m_u32SectorsPerBlock;</pre>	//!< Number of 2K sections per block.	
u	int32_t	m_u32NumberOfNANDs;	//!< Total Number of NANDs - not used by ROM.	
u	int32_t	m_u32TotalInternalDie;	//!< Number of separate chips in this NAND.	
u	int32_t	m_u32CellType;	//!< MLC or SLC.	
u	int32_t	<pre>m_u32EccBlockNEccType;</pre>	//!< Type of ECC, can be one of BCH-0-20	
u	int32_t	m_u32EccBlock0Size;	//!< Number of bytes for Block0 - BCH	
u	int32_t	<pre>m_u32EccBlockNSize;</pre>	//!< Block size in bytes for all blocks other than Block0 - BCH	
u	iint32_t	<pre>m_u32EccBlock0EccType;</pre>	//!< Ecc level for Block 0 - BCH	

IMX UBOOT GITHUB REPOSITORY

int ecc_strength;

{

```
/*
 * Determine the ECC layout with the formula:
        ECC bits per chunk = (total page spare data bits) /
 *
 *
                (bits per ECC level) / (chunks per page)
 * where:
 *
        total page spare data bits =
 *
                (page oob size - meta data size) * (bits per byte)
 */
ecc_strength = ((page_oob_size - MXS_NAND_METADATA_SIZE) * 8)
                / (MXS NAND BITS PER ECC LEVEL *
                        mx28_nand_ecc_chunk_cnt(page_data_size));
```

Y U NO ENCRYPT ?

- i.MX systems support NAND flash encryption
- Most of the systems we have tested so far do not use encryption (what did you expect ?)

KNOWN VARIANTS

- Some i.MX dumps we made seemed to use a different ECC mechanism
- Various GPMI drivers mention **different versions** of Freescale ROM and variants of FCB structure
- The current version of *imx-nand-tools* worked for all of our dumps but may fail with yours, so ... digital.security

INSTALL, TEST, AND CONTRIBUTE !

CONCLUSION

- i.MX system uses a **custom NAND flash layout**
- This **layout is documented** in various documents and publicly available code
- **imx-nand-tools** provides a set of tools to handle this layout and convert dumps into useable images
- i.MX systems **should use NAND flash encryption feature** to avoid key/password/IP leaks

THANKS FOR ATTENDING, ANY QUESTION?

Contact Additional Add

RELATED LINKS

- PTP firmware extraction tips & tricks:
 - https://www.pentestpartners.com/security-blog/how firmware-analysis-tools-tips-and-tricks/
- IMX28 Reference manual: https://bootlin.com/~maxime/pub/datasheet/MCIMX2
- UBOOT NAND utility: https://github.com/u-boot/uboot/blob/master/tools/mxsboot.c
- Freescale Linux driver: https://github.com/Freescale/ fslc/tree/4.1-2.0.x-imx/drivers/mtd/nand/gpmi-nand