

SECURITY VULNERABILITIES OF DIGITAL VIDEO BROADCAST CHIPSETS

Adam Gowdiak Security Explorations

INTRODUCTION



About Security Explorations

- Security start-up company from Poland
- Provides various services in the area of security and vulnerability research
- Commercial and Pro Bono research projects
- Came to life in a result of a true passion of its founder for breaking security of things and analyzing software for security defects
- Our ambition is to conduct quality, unbiased, vendor-free and independent security and vulnerability research

INTRODUCTION



Presentation Goal

- Continuation of our research in a digital satellite TV area
- Educate about security risks associated with less known technologies and platforms such as those used in a digital satellite TV ecosystem
- Warn about security risks associated with
 - closed ecosystems such as digital satellite TV
 - insecurely implemented proprietary hardware components
 - 3rd party security evaluation processes

INTRODUCTION

DISCLAIMER



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http://www.security-explorations.com/en/legal-threats.html



Why bother about content security ?

- Pay TV piracy remains a major concern for channels and operators
 - it leads to financial losses for the European pay TV industry
 - it substantially damages the image of transmitters and content rights holders
 - it reduces the allure and payback of investing in the industry
 - it hurts the industry and its innovation capabilities
- Signal theft estimated to be more than \$2.1 billion at the end of 2011 for Asia region alone (CASBAA)

Security of a premium content



- Paid, premium content broadcasted in encrypted form
 - Scrambling at the TS or PES level
 - transport_scrambling_control bit of MPEG TS packet
 - Common Scrambling Algorithm (CSA) and its derivatives
 - Shared 64-bit secret key (Control Word)
 - Dedicated security chipsets for decryption
- Key components in the security system
 - Subscriber's smartcard
 - holds information about subscriber's access rights to programming
 - releases decryption keys to the set-top-box if access to a given service is granted
 - Set-top-box
 - Conducts decryption of a scrambled content with the use of a received decryption key

Security of a premium content (2)







Control Words (CW)

- 64-bit secret keys used to descramble encrypted MPEG streams
 - Audio, video and data
- Unique to each programming
- □ Generated automatically by the content provider
 - Changed every ~10s
 - Odd and even keys for uninterrupted programming reception
 - Current and next key
- Broadcasted in encrypted form to client devices (set-top-boxes)
 - carried in entitlement control messages (ECM)
 - encrypted with the use of asymmetric crypto (i.e. RSA)



Entitlement Control Messages (ECM)

- ECM messages contain private conditional access information such as Control Words
 - Broadcasted by the means of a dedicated MPEG stream
 - Message format specific to CAS vendor
- PID of MPEG stream carrying ECM messages denoted by CA_descriptor
 - If elementary stream is scrambled, a CA descriptor shall be present for the program containing that elementary stream
 - Usually present in TS_program_map_section
 - MPEG table_id = 0x02



Conditional Access System (CAS)

- It protects the content by requiring certain criteria to be met before granting access to the content
- Subscriber's smartcard holds information about subscriber's access rights to a given programming
 - what programming / program packages a subscriber is entitled to watch
- Only authorized client devices (paying subscribers) can decrypt MPEG streams for premium content
 - Set-top-box device asks the smartcard to decrypt encrypted Control Word (ECM message)
 - The smartcard makes sure that access to the content can be granted and releases the plaintext value of a Control Word

CAS architecture (set-top-box side)







Threats to the model

- Premium content is encrypted and broadcasted to all subscribers with the use of same crypto key (Control Word)
- One rogue subscriber with access to all premium content can share Control Word keys with others over the Internet
 - illegal reception / distribution of premium programming aka signal theft
 - Control Words sharing

security explorations

CAS with chipset pairing

- Control Words unique for each pair of a subscriber (smartcard) / client device (set-top-box)
 - smart cards can be used only with secure devices
 - the link between the smart card and the client device is secured
 - illegal content redistribution is prevented (no more CW sharing)
- Chipset pairing has a form of a cryptographic function
 It is usually implemented in a silicon chip (DVB chipset)







Pairing function



A function that cryptographically ties a set-top-box device and a subscriber's smartcard
 encCW = PAIRING_{enc}(CW,CWPK)
 CW = PAIRING_{dec}(encCW,CWPK)

Control Words pairing key (CWPK)

- Unique to each subscriber
- Assigned to it at the time of activating a given user's digital satellite TV subscription
- Usually, a function of a unique DVB chipset's key

Conax CAS with chipset pairing



- Conax AS is one of the major CAS providers for the Pay TV industry
 - More than 350 installations in 80 countries world-wide
- CAS implemented in software and hardware
 - Partnership with set-top-box vendors
 - Partnership with many DVB chipset vendors to implement chipset pairing functionality
 - STMicroelectronics, Broadcom, Renesas Electronics, ...





- Mandatory, comprehensive evaluation of all qualified chipsets run at independent, world-leading security laboratories
 - All set-top-boxes and DVB chipsets implementing Conax CAS with chipset pairing undergo rigorous security evaluation process
 - Official scoring assigned to set-top-boxes and DVB chipsets and certified in writing by Conax CSO
 - "O represents no security and 9 corresponds to the security level of Conax smart cards"

Source: http://www.conax.com/en/solutions/clientdevicesecurity/ Conax Security Department (09-Jan-2012)



Conax security certification



Source: Neotion company website

Introduction



- DVB chipsets implement the core functionality related to the handling of MPEG transport streams and A/V data such as:
 - MPEG transport filtering and descrambling (incl. chipset pairing function)
 - audio and video decoding
 - graphics display
 - communication interfaces
 - memory interfaces
- For security and efficiency reasons, they are usually implemented as a single chip (system-on-chip or SoC)
 - Multiple processor cores for various functions



STMicroelectronics implementation

STi7100 single-chip, high-definition STB decoder



source: st.com



STMicroelectronics implementation (2)

□ STi7111 single-chip, high-definition STB decoder



source: st.com



STMicroelectronics STB H.264 generations



Source: Multimedia Convergence & ACCI Sector Overview, Philippe Lambinet, STMicroelectronics



Security challenges

- Implementation of a chipset pairing function in a proprietary silicon chip makes it far more difficult to reverse engineer and break
 - no target software for the static analysis / reverse engineering or runtime interception
 - undocumented interfaces
 - unknown implementation of the pairing function
 - unknown crypto algorithm and keys (their sizes, byte order, etc.)





- Tedious analytical and reverse-engineering work
- By gathering and gluing together many pieces of information (clues), it was possible not only to discover the operation and implementation of investigated DVB chips, but also find security weaknesses in them
- \Box The tools
 - Without custom reverse engineering tools we would not be able to successfully complete most of our projects
 - This is especially valid for SE-2011-01 project

Common approach (chips documentation)



Data briefs available from st.com (STi710x, STi7111, STM7710, etc.)

- Generic chip architectures
- Processor cores
 - ST40 32-bit superscalar RISC CPU
 - Dual ST231 CPU cores for audio and video decoding
- Transport subsystem
 - Programmable Transport Interface(PTI)
 - PID filtering, Demultpilexing, Descrambling
 - Transport Stream Merger (TSM) and router
- FDMA controller
 - PES parsing and start code detection
 - Routing elementary streams to A/V buffers



Common approach (discovering core device drivers)

- Device drivers implementing Control Words operations
- Static / binary analysis
 - Inspecting libraries and device driver code / symbols
 - Figuring out code dependencies
 - Call and link graphs

GSECHAL device driver (STi7100)

gSecHAL_Init gSecHAL_GetRevision gSecHAL_SetAlgorithm gSecHAL_CKCalc gSecHAL_DecryptSCK gSecHAL_GetStatus gSecHAL_CopyTCN gSecHAL_Reset

STTKDMA device driver (STi7111)

STTKDMA_Reset STTKDMA_DecryptKey STTKDMA_ReadPublicID sttkdmaHal_GetNonce STTKDMA_GetCounter STTKDMA_Nop sttkdmaHal_GetSWReg

. . .



Common approach (pinning down CW API calls)





Helpful CAS system implementation

- Non-HD prepaid satellite TV service (TNK) available along the main Platform 'N'
 - Different set-top-box decoders
 - Technisat, ...
 - Conax CAS smartcards
- A few services available to both SAT TV platforms
 - Shared audio / video streams (same broadcast)
 - Separate conditional access information
 - separate ECM streams



Helpful CAS system implementation (2)

Program Stream Infromation



Helpful CAS system implementation (3)

security explorations

- Parallel Conax CAS without chipset pairing
- The plaintext values of encrypted Control Words





Device drivers' functionality and operation

Detailed analysis of GSECHAL's DecryptSCK function

- The meaning of configuration data
 - SecureMode = 1
 - UsingAES = 0
- Memory mapped I/O registers
 - CONFIG
 - STATUS
 - COMMAND
 - DATA
- Implementation of direct chip programming commands
 - gSecWaitForComplete, gSecDataRead, gSecDataWrite,

Reconstruction of a chip's programming sequences



DecryptSCK IOCTL (pseudocode) If (SecureMode=1, UsingAES=1) { gSecDataWrite buf 4 gSecInstWrite 3|arg<<8 If (SecureMode=0, UsingAES=1) { d1 -> BASE_ADDR_+0x6100+arg<<4 d2 -> BASE ADDR +0x6104+arg<<4 d3 -> BASE ADDR +0x6108+arg<<4 d4 -> BASE_ADDR_+0x610c+arg<<4 If (SecureMode=1, UsingAES=0) { gSecDataWrite buf 2 gSecInstWrite 3|arg<<8 gSecInstWrite 3|arg<<8 gSecSetKeyPtr If (SecureMode=0, UsingAES=0) { d1 -> BASE ADDR +0x6100+arg<<4 d2 -> BASE ADDR +0x6104+arg<<4 d3 -> BASE_ADDR_+0x6108+arg<<4 d4 -> BASE ADDR +0x610c+arg<<4 gSecSetKeyPtr

DecryptSCK implementation leaks Control Words storage addr (0x6100)

- Default GSEC chip operation
- SecureMode value defined by chip fuses



GSEC keys memory

- Secure storage area for Control Word keys
 - Ox6100 offset from the chip's base addr
 - Unavailable for reading / writing
 - Read operation always returns ZEROs
 - Write operation does not disrupt the descrambling process
- The arguments to the DecryptSCK command include the index of the key slot to load with 0x10 bytes (key data)
 - Device driver code makes sure that this index is within the 0x00-0x31 bounds



Security vulnerability

- The ability to extract plaintext values of Control Words
 - the chip needs to be programmed manually by issuing commands directly to its I/O mapped registers
 - index of the DecryptSCK command needs to be greater than 0x31

plaintext CW mem	=	0x6100 + key_idx * 0x10
		where key_idx>0x31



Security vulnerability (formula)

$Key_{32h} \leftarrow TDES_{dec}(encCW, plainCWPK)$ $Key_{32h} \equiv plainCW$

Where:

Key_{32h} = key slot #32h (accessible chip location) encCW = encrypted Control Word value (known value) plainCWPK = plaintext Pairing Key value (unknown value)



Security vulnerability (SoC location)



original image: st.com
security explorations

Introduction

- Different design brings more challenges
 - STTKDMA chip component
 - One DecryptKey IOCTL for key related operations
 - Control Words
 - Pairing Key (CWPK)
 - AES keys
- □ Easier reverse engineering
 - Modular architecture of ITI2850ST/ ITI2849ST set-top-boxes' OS distribution
 - Many dedicated user level libraries
 - Conax CA, Crypto operations, NAND encryption, STB configuration, ...
 - Text XML configuration files
 - Conax CA client settings
 - Support for crypto DMA operations
 - Kernel symbols via /proc/kallsyms (680KB+)



Dedicated key memories

- Separate memory mapped chip regions for AES and Control Word keys
 - Deduced with the help of kernel and user level library symbols
 - Code / data symbols associated with accesses to chip's I/O memory

DecryptKEY FUNCTIONALITY

cpcw_keys, descrambling keys (set_cleartext_descramblingkey) 3100 offset, indices 0x00-0x31, key size 0x10

CRYPTO DMA FUNCTIONALITY

sttkdma keys, cdma_dev_keys, cpcwKeys (STDRMCRYPTO_AES_LoadKeySlot) 3420 offset, indices 0x00-0x07, key size 0x10

BREAKING STi7111 CHIPSET CCORE library



- The library used for crypto DMA operations
- Reverse engineering the meaning of CCORE library/ chips configuration bits and input / output arguments
 - Manual analysis of data propagation
 - Ibstd_drv_ccore.so API -> STTKDMA device driver API -> chip's configuration registers
 - DMA CONFIG
 - TKD CONFIG
 - Analysis of STTKDMA code writing to configuration registers
 - resetAES_NOT_TDES
- Custom AES / TDES Java subroutines
 - Verification of CCORE results



The existence of a chip specific SCK key

- Initial hints in STi7100 GSECHAL device driver
 - DecryptSCK command
- Confirmed at a time of the analysis of a set-top-box boot loader code (SH4 emulator)
 - The use of SCK key to decrypt the boot loader code
 - No initialization of the usual key registers
 - Different chip configuration bits
- Used by a device driver from a software upgrade OS distribution
 - The use of SCK key for NAND encryption
 - parm=nand_crypt_use_sck_key:Use SCK key instead of default ADB key for NAND encryption



Firmware data / code for STi7111 chip

- Hints in sttkdma_core_user.ko module
 - Code symbols
 - st_tkdma_loader
 - st_tkdma_loader_checksum
 - Data symbols
 - tkdma_fw_address_1
 - tkdma_fw_address_2
 - Writing data / code to STi7111 chip's I/O space
 - 5944 code bytes
 - 1156 data bytes
- Firmware code implementing unknown processor instructions



SLIM Core processor

- IST FP6 PROSYD EU project (<u>http://www.prosyd.org</u>)
- Paragraph 2.3 of Deliverable D1.4/1 gives some information about the SLIM Core Processor
 - a collaboration between ST UK and OneSpin after the spin-off from Infineon
 - Iightweight processor with 27 instructions and a 4-stage pipeline
 - processor special features: a coprocessor interface; circular buffer operation; a STOP instruction
 - Instructions opcode names: ADD BRA CPI JAB LD LDF NOP RPT STI STF STOP SUBC.



JMP instruction format

OpenDuckbox project (<u>http://gitorious.org/open-duckbox-project-sh4</u>)

- GNU source code for SLIM Core Generic driver
- slim_boot_core function leaks information about the format of one Slim Core instruction (JMP)
 - memory addressing (by word number)
 - instruction opcode width



Finding patterns in SLIM Core code



Result of GetPublicID command

Discovery of STORE instruction format

- Firmware pattern matching format of GetPublicID result
- Confirmed by changing ST R5 with ST R0 instruction



Reverse engineering SLIM Core instructions

- Discovery of SLIM Core instruction opcodes
 - Exploited the ability to change the operation of SLIM
 Core firmware in runtime
 - Overwriting chip's memory loaded with firmware code
 - should secure crypto chip allow for it ?
 - Replacement of an arbitrary instruction from the code path of GetPublicID function
 - Analysis of the instruction's execution effect to memory and registers
 - Discovery of load / store instructions format



Reverse engineering SLIM Core instructions (2)



Reverse engineering SLIM Core instructions (3)



- JMP and LOAD/STORE instructions sufficient to discover the meaning of all other instructions
 - JMP from firmware to user's code path
 - STORE the contents of registers (firmware context)
 - LOAD user's environment (contents of registers)
 - **EXECUTE** unknown SLIM Core instruction opcode
 - **STORE** user's environment (contents of registers)
 - LOAD the contents of registers (firmware context)
 - JMP back to firmware code path
- The need to properly handle conditional jumps



Reverse engineering SLIM Core instructions (4)

- One instruction opcode at a time
 - LOAD / STORE instructions
 - MOV instructions
 - CMP instructions
 - Conditional branching instructions
 - Computational instructions
 - Other instructions (bit extraction, manipulation)
 - RPT instruction
- Scope limited to unknown opcodes from firmware code

Reverse engineering SLIM Core instructions (5)



Visible code patterns

MOV instruction patterns



MOV instruction confirmed by the result of STORE

CMP / conditional jump instruction patterns

0014	0x <mark>00981</mark> 026	je 1_0026
0015	0x00c030 <mark>02</mark>	cmp r3, #02
0016	0x00981028	je 1_0028
0017	0x00c030 <mark>06</mark>	cmp r3, #06
0018	0x <mark>00981</mark> 02a	je 1_002a
0019	0x00c0300b	cmp r3, #0b
001a	0x0098102c	je 1_002c
001b	0x00c030 <mark>0f</mark>	cmp r3,#0f
001c	0x0098102e	je 1_002e
001d	0x00c030 <mark>03</mark>	cmp r3, #03
001e	0x00981030	je 1_0030
001f	0x00c030 <mark>07</mark>	cmp r3,#07

CMP values correspond to device driver commands



SLIM Core disassembler

Final disassembly dump of Slim Core firmware code

- 1400+ instructions disassembled
- ~11 instruction opcodes not recognized
 - Not relevant from the analysis point of view
- Sufficient data for firmware analysis
 - Discovery of separate dispatching for DMA and all TKD operations
 - Semi-threads (context-switching)
 - Discovery of a key initialization subroutine



Tracing SLIM Core firmware

- The goal was to locate SLIM Core instruction sequences implementing DecryptKey functionality
- Tracer implementation
 - SLIM core part
 - Custom code on the GetPublicID function path
 - Binary instrumented instruction copied from a currently traced code location
 - SLIM Core instruction executed in the original registers context
 - Heavy use of the SLIM core disassembler
 - Java part
 - Logging
 - SLIM Core syncing and control code



Tracing SLIM Core firmware (output log)

starting logger at 80 break at: 0x0000086 r0 00000000 *r1 00000001 *r2 00000100 r3 00000000 *r4 00000011 *r5 00000031 *r6 00001103 *r7 00000005 *r8 00000006 *r9 31ff0001 r10 00000000 *r11 00000001 r12 00000000 *r13 0000024e *r14 000000d0 IP 00000086 0086 0x00e10001 mov r1,#0001 SCDC stopped mov r1,#0001 0086 0x00e10001 break at: 0x0000087 r0 0000000 r1 0000001 r2 0000100 r3 0000000 r4 00000011 r5 00000031 r6 00001103 r7 00000005 r8 0000006 r9 31ff0001 r10 0000000 r11 00000001 r12 0000000 r13 000024e r14 000000d0 IP 0000087 0087 0x00a20048 ld r2,[r0,0048] // 0x4120 break at: 0x0000088 r0 0000000 r1 00000001 *r2 00000001 r3 00000000 r4 00000011 r5 00000031 r6 00001103 r7 00000005 r8 0000006 r9 31ff0001 r10 0000000 r11 00000001 r12 00000000 r13 0000024e r14 000000d0 IP 00000088 0088 0x00722c21 bitval r2,r2,#0002 0x00881091 0089 jz 1 0091 break at: 0x00000091 r000000000 r1 00000001 *r2 00000000 r3 00000000 r4 00000011 r5 00000031 r6 00001103 r7 00000005 00000006 r9 31ff0001 r10 00000000 r11 00000001 r8r12 00000000 r13 0000024e r14 000000d0 IP 00000091

Tracing SLIM Core firmware (DecryptKey code)

DecryptKey implementation

			must alsigning TDE0 mode
0206 0x00fa4000	COPTDES		put chip into IDES mode
0207 0x000f093c	mov r15,r9		load TKD command
0208 0x008e1208	wait1		wait
0209 0x00af0008	ld r15,[r0,0008]		load encrypted CW
020a 0x00af0009	ld r15,[r0,0009]		
020b 0x00af000a	ld r15,[r0,000a]		
020c 0x00af000b	ld r15,[r0,000b]		
020d 0x008e120d	wait1		wait
020e 0x00500a00	tst r10,r10		
020f 0x00881215	jz 1_0215		
0210 0x00b0f008	st r15,[r0,0008]		optionally store result
0211 0x00b0f009	st r15,[r0,0009]		
0212 0x00b0f00a	st r15,[r0,000a]		
0213 0x00b0f00b	st r15,[r0,000b]		
0214 0x00d02117	jmp 1_0217		
0215 0x00d00004	rpt 4		
0216 0x00000f3c	mov r0,r15		
0217 0x00d0211a	jmp 1 021a		
		,	





Internal crypto core (TKD) commands

Static analysis of SLIM Core firmware disassembly

Discovery of internal chip commands

001181	stk cr	nd 2	setCWPK	
01ff81	101 stk cr	nd 1		
02ff81	101 stk cr	nd 0x20		
03ff00	001 stk cr	nd 0x10		
10ff80	001 stk cr	nd idx<<8	0x80	
10ff01	101 stk cr	nd idx<<8	0x04	
20110	01 stk cr	nd idx<<8	0x03 set_protected_descramblingke	Y
20ff00	010 stk cr	nd idx<<8	0x06	-
040000	001 stk cr	nd 0x11		0.00//0404 0110
			Set CWPK	0X00ff8101 CMD
ffff04	401 stk cr	nd $0x12$	Set encrypted CW idx 0x00	0x20ff0001 CMD
80ff02	203 stk cr	nd 0x21	Set encrypted CW idx 0x01	0x21ff0001 CMD
81ff02	203 stk cr	nd 0x22	Set encrypted CW idx 0x02	0x22ff0001 CMD
82ff02	203 stk cr	nd 0x23	Sot onorupted CW idx 0x02	0x22ff0004 CMD
83ff02	203 stk cr	nd $0x24$		
83ff02	203 stk cr	nd 0x24		
83ff02	203 stk cm	nd 0x24	Set encrypted CW idx 0x31	0x51ff0001 CMD

TKD inspector



Discovery of the meaning of TKD commands by the means of executing special SLIM Core instruction sequence

public static final int tkd code[] = {		
0x00e61234,	//mov r6,#1234 TKD_CMD_HI	TKD command to test
0x00e55678,	//mov r5,#5678 TKD_CMD_LO	
0x00e00000,	//mov r0,#0000	
0x00756210,	//mov r5,(r6<<16) r5	
0x00e00000,	//mov r0,#0000	
0x00fa4000,	//COPINS	Chip mode instruction
0x00e00000,	//mov r0,#0000	(AES/TDES)
0x000f053c,	//mov r15,r5	
0x00e00000,	//mov r0,#0000	Wait ins (different for
0x008e1abc,	//WAITINS	AES/TDES mode)
0x00e00000,	//mov r0,#0000	Input data (kov)
0x00af0050,	//ld r15,[r0,0050]	input data (key)
0x00af0051,	//ld r15,[r0,0051]	
0x00af0052,	//ld r15,[r0,0052]	
0x00af0053,	//ld r15,[r0,0053]	
0x00e00000,	//mov r0,#0000	Wait ing (different for
0x008e1abc,	//WAITINS	Wait ins (different for
0x00e00000,	//mov r0,#0000	AES/IDES mode)
0x00b0f054,	//st r15,[r0,0054]	Output data (key)
0x00b0f055,	//st r15,[r0,0055]	
0x00b0f056,	//st r15,[r0,0056]	
0x00b0f057,	//st r15,[r0,0057]	
0x00e00000	//mov r0,#0000	-
};		



TKD inspector (playing with the commands)

```
test> input "b6 04 78 c3 0f 26 a3 06 d5 20 10 0f c0 93 4f f3"
test> ed 0x01ff0000 0x00fa4000 0x008e1abc
tkcmd 01ff0000
[running SLIM code]
test> input "e0 c9 cd 4e 2f bd 52 a0 e0 c9 cd 4e 2f bd 52 a0"
test> ed 0x15ff0101 0x00fa4000 0x008e1abc
tkcmd 15ff0101
[running SLIM code]
test> kevs
[00] 6d 31 99 ca f8 fa e5 51 fc 56 22 11 c2 33 05 8a
[01] 6d 31 99 ca f8 fa e5 51 fc 56 22 11 c2 33 05 8a
[02] 6d 31 99 ca f8 fa e5 51 fc 56 22 11 c2 33 05 8a
[03] f8 25 41 20 00 00 00 00 00 00 00 00 00 43 11 71
[05] 00 85 55 26 86 09 29 54 00 85 55 26 86 09 29 54
[07] d4 c8 94 af 84 84 5c de 17 82 f7 73 1e c3 2f e7
DMA CONFIG: 20 08 00 00
TKD CONFIG: 00 00 00 00
INPUT: e0 c9 cd 4e 2f bd 52 a0 e0 c9 cd 4e 2f bd 52 a0
```



TKD commands format



TARGET:

- target, where the result of the operation should be stored
- a key slot number or 0xff for chip registers

SOURCE:

- source, from which data for the operation should be fetched
- a key slot number or 0xff for chip registers,

KEY:

- key slot number, which holds the key used for the crypto operation
- value 0x00 usually identifies SCK key (unique key for each chip)

CONFIG:

- configuration bits
- bit 0 Usually denotes encryption (0) or decryption (1) operation



TKD commands format (key operations explained)

TKD CMD 0x00ff0000

- Setting encrypted Control Word Pairing Key (CWPK)
- Interpreted as decryption (always) of register input (0xff) with SCK key (0x00) and storing the result at a key slot 0x00

TKD CMD 0x20ff0001

- Setting encrypted Control Word
- Interpreted as decryption (0x01) of register input (0xff) with SCK key (0x00) and storing the result at a key slot 0x20

Security vulnerability #1 (step 1)







Security vulnerability #1 (step 2)





Security vulnerability #1 (formula)

 $Key_1 \leftarrow TDES_{dec}(encCWPK,SCK)$ $Key_{15h} \leftarrow TDES_{dec}(encCW,Key_1)$ $Key_1 \equiv plainCWPK$ $Key_{15h} \equiv plainCW$

Where:

Key₁ = key slot #01h (inaccessible chip location) Key_{15h} = key slot #15h (accessible chip location) encCW = encrypted Control Word value (known value) encCWPK = encrypted Control Word Pairing Key value (known value) SCK = chip specific key (unknown value)



Security vulnerability #2 (step 1)





Security vulnerability #2 (step 2)





Security vulnerability #2 (formula)

$Key_{15h} \leftarrow TDES_{dec}(plainCWPK,SCK)$ $plainCWPK \leftarrow TDES_{enc}(Key_{15h},SCK)$

Where:

Key_{15h} = key slot #15h (accessible chip location) SCK = chip specific key (unknown value) plainCWPK = plaintext Pairing Key value (unknown value)



STi7111 security vulnerabilities (SoC location)



original image: st.com



Control Word Pairing Key

- Issue 18 makes use of the encrypted value of the chipset pairing key (CWPK)
- CWPK key sent by the operator at the time of activating user's subscription
 - Encrypted CWPK key bytes returned by the Conax card in response to EMM message
- For ITI2850ST and ITI2849ST set-top-boxes encrypted CWPK key encrypted again and cached in a file
 - /mnt/flash/secure/7/0
 - Ibstd_cai_client_conax7.so API for decryption

PROOF OF CONCEPT CODE



Implemented functionality

- Access to information about cryptographic keys
 - Plaintext Control Words (STi7100)
 - Plaintext Control Words and plaintext CWPK (STi7111)
- Control Words sharing via network between arbitrary decoders protected with Conax conditional access method and chipset pairing
- Video on Demand ECM decryption and sharing of programming protected with Conax conditional access method with chipset pairing

PROOF OF CONCEPT CODE



Control Words sharing (aka. signal theft)

Sharing of the crypto keys used to descramble digital satellite
TV programming



PROOF OF CONCEPT CODE



Push VOD sharing

- Sharing of the crypto keys used to descramble VOD movies
 - Obtaining Control Word for arbitrary movie during the rental period
 - VOD movies rented for 48 hours period
 - Encrypted MPEG data pushed into set-top-boxes
 - ECM messages accompanying movies files
 - Sharing Control Words after the rental period
 - Using plaintext CW values to descramble the movie
 - The use of key memory beyond index 0x32 on STi7100
 - The use of CWPK to reencrypt Control Word on STI7111

SUMMARY



STi7100 / STi7111 security vulnerabilities

- How come the issues were not discovered before the market release ?
 - STMicroelectronics a major silicon vendor
 - #1 in Europe , #7 in the world (source: Wikipedia)
 - Conax security evaluation of STB / CAM / DVB chipset solutions
 - Final scoring of STi7100/STI7111 not disclosed to Security Explorations





- **Vulnerabilities Impact**
- No information from STMicroelectronics (DVB chipsets vendor) in response to the impact inquiry questions
 - All your inquiries, as listed below, are pointing towards confidential information and as such can not be disclosed by ST to you or to others.

Jan-17-2011, STMicroelectronics in an email to Security Explorations

Impact estimation upon publicly available data





Vulnerabilities Impact (2)

Cumulative MPEG-2 & MPEG-4 Shipments in 2008

541 millions of units

- Set-top-boxes, digital television sets, DVD / Bluray players
- STMicroelectronics #1 in H.264 market (68% of market share in 2008)
- Customers from Europe, Middle East and Africa, Asia-Pacific and the Americas
 - DishTV (India)
 - DirectTV (USA)
 - Platforma N, Cyfrowy Polsat (Poland)
 - BSkyB (UK)
 - ••••

Source: Multimedia, Philippe Lambinet, STMicroelectronics
SUMMARY

Final Words



- First successful attack against the implementation of a Conax conditional access system with chipset pairing
 - Pay TV piracy possible in the environment of hacked digital satellite TV set-top-boxes
 - Security of dedicated DVB chipsets broken
- Security based on a complex, secret functionality embedded in a silicon is a dangerous concept

Security through Obscurity ?

- The need to improve security and evaluation processes by silicon and CAS vendors
- The need to tighten set-top-boxes security relying on vulnerable DVB chipsets

HITBSecConf, May 24-25, 2012, Amsterdam, The Netherlands



Q & A



THANK YOU

contact@security-explorations.com

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